STL Gotchas

Avoiding Common Errors in Using the STL

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- Insertion and deletion might invalidate references, pointers, and iterators.
- A set iterator must not allow modification of the elements.
- Function objects must neither have side effects nor modify container elements.
- Comparators must not be polymorphic.
- Associative containers use induced equivalence, whereas algorithms use equality.
- Adapted iterators cannot be passed to container operations.
- Stream iterators on the same stream are not independent of each other.
- Allocators must exhibit static behavior.

STL Pitfall #1

Invalidation of references, pointers, and iterators

All container in the STL ...

- are of dynamic size, i.e. they grow as needed
- allow insertion and removal of elements via member function i nsert() and erase()
- provide iterators that give access to the contained elements
- provide iterators to the beginning and end of the sequence

Consider a program that reads lines from an input file, sorts them, and writes the result of sorting to an output file, using a l i st as the temporary store.

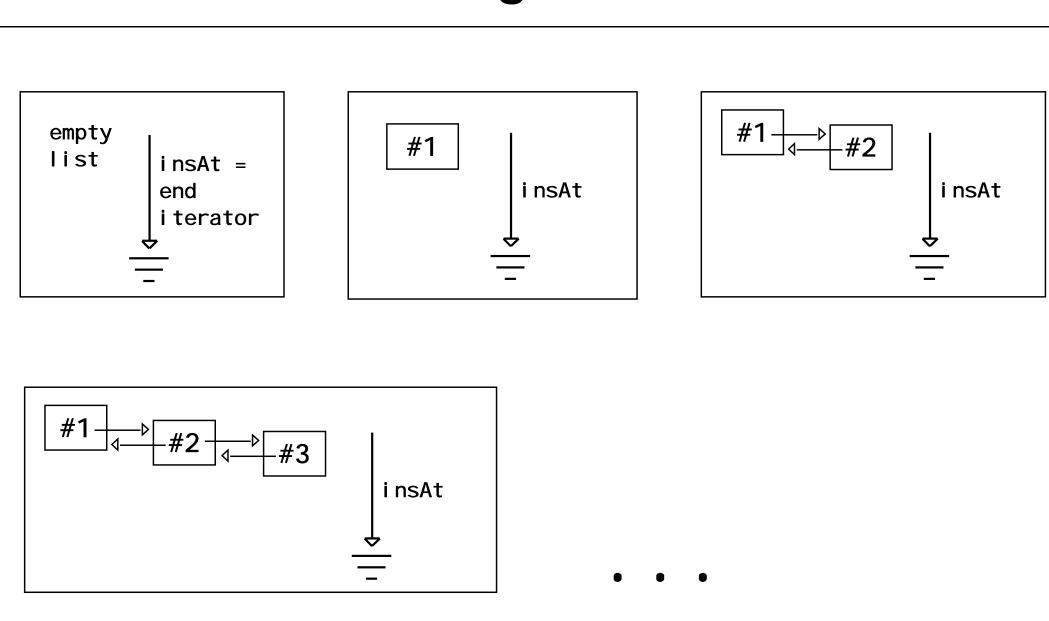
Inserting elements to a list

```
void dolt(const char* in, const char* out)
{ list<string> buf;
   list<string>::iterator insAt = buf.end();
   string linBuf;
   ifstream inFile(in);
   while(getline(inFile,linBuf))
          buf.insert(insAt,linBuf);
   buf.sort();
   ofstream outFile(out);
   copy(buf.begin(), buf.end(),
        ostream_i terator<string>(outFile, "\n"));
```

Insert at a specified position

Inserts a copy of value before the specified position.The returned iterator points to the newly inserted copy of value.

Inserting elements to a list



```
void dolt(const char* in, const char* out)
{
  set<string> buf;
   set<string>::iterator insAt = buf.end();
   string linBuf;
   ifstream inFile(in);
   while(getline(inFile,linBuf))
          buf.insert(insAt,linBuf);
   ofstream outFile(out);
   copy(buf.begin(), buf.end(),
        ostream_i terator<string>(outFile, "\n"));
```

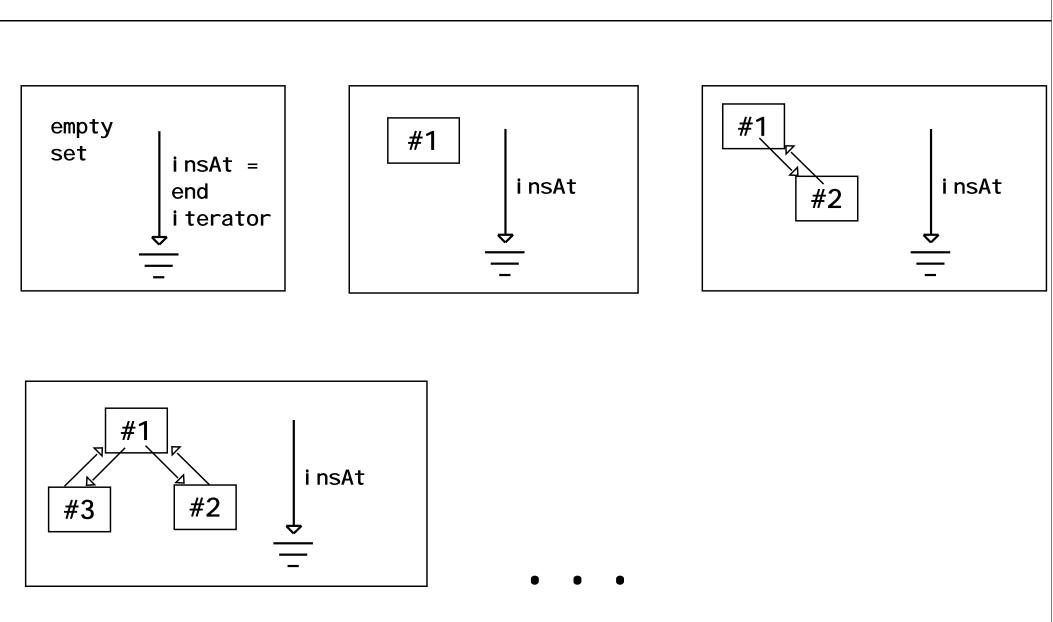
Insert with hint

Inserts a copy of value.

- Inserts only if there is no element in the container with the same value.
- SThe returned iterator points to the element with the same value.

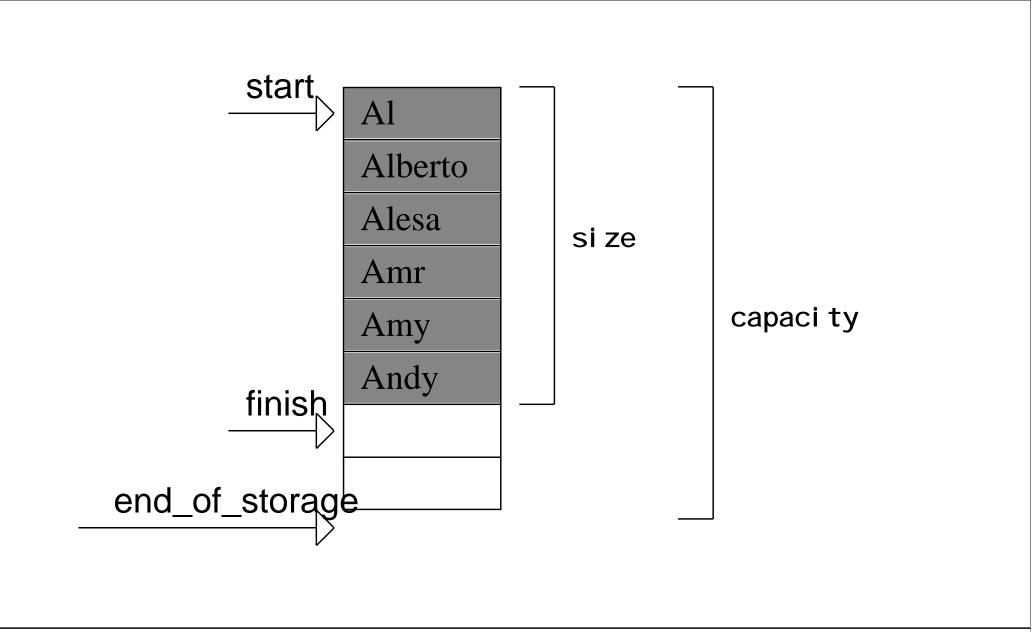
SThe iterator position is a hint. It should point to where the value should be inserted.

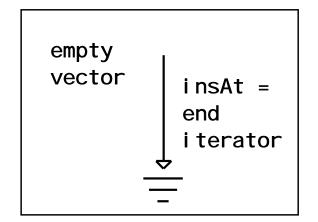
A The hint has no effect on correctness, only on performance.

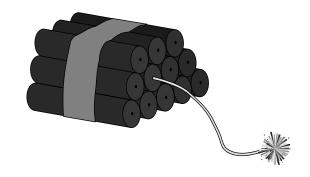


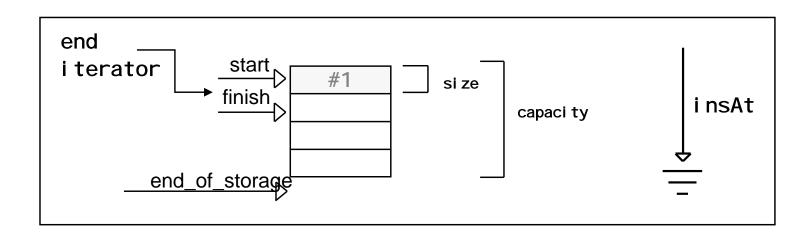
```
void dolt(const char* in, const char* out)
{
  vector<string> buf;
   vector<string>::iterator insAt = buf.end();
   string linBuf;
   ifstream inFile(in);
   while(getline(inFile,linBuf))
                                    <<<< crash !!!</pre>
         buf.insert(insAt,linBuf);
   sort(buf.begin(), buf.end());
   ofstream outFile(out);
   copy(buf.begin(), buf.end(),
        ostream_i terator<string>(outFile, "\n"));
```

Typical implementation of vector









So, what's going wrong ... ?

- dol t () is built on the assumption that an iterator (i nsAt = buf. end()), that is valid in one context (with an empty vector), is still valid in another context (after insertion to the vector).
- There is no such guarantee.

```
void dolt(const char* in, const char* out)
{ ...
   vector<string>::iterator insAt = buf.end();
   ...
   while(getline(inFile,linBuf))
        buf.insert(insAt,linBuf);
   ...
}
```

Validity guarantees for insertion

For associative containers and list:

• Does not affect the validity of iterators and references.

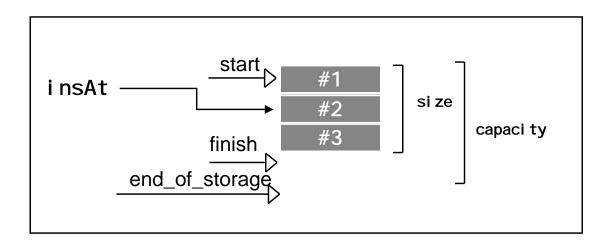
For vector:

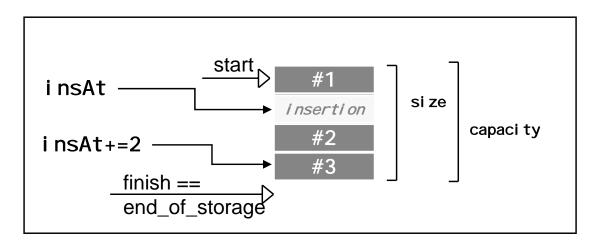
- Causes reallocation if the new size is greater than the old capacity.
- Reallocation invalidates *all* the references, pointers, and iterators.
- If no reallocation happens, all the iterators and references before the insertion point remain valid.

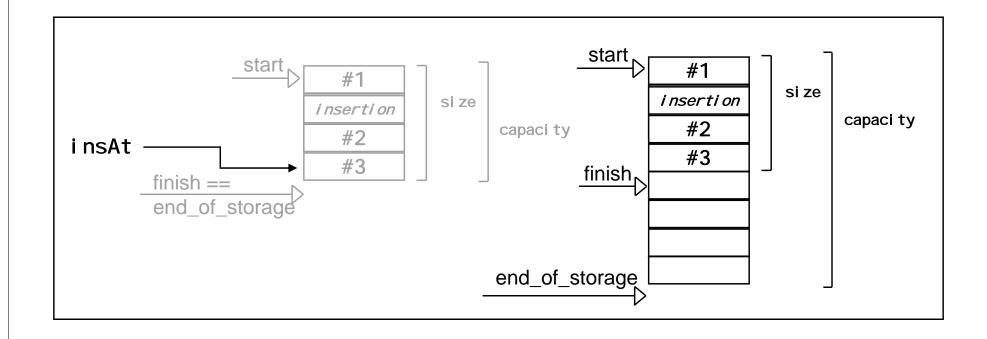
- Does not only apply to the end iterator, but to *any* iterator after the point of insertion:
- Insertion moves all elements after the point of insertion to the back.
- All references to elements after the point of insertion become invalid.
- In particular, the point of insertion itself becomes invalid as a side effect of the insertion.

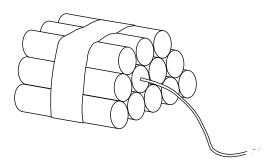
Consider a program that inserts a line at a specific position, say, before any line that starts with a capital letter.

```
void dolt(const char* in, const char* out)
{
 vector<string> buf;
   // ... populate vector ...
   vector<string>::iterator insAt = buf.begin();
   while ((insAt = find_if(insAt, buf.end(), isUpper()))
          ! = buf. end())
                                      <<<<< crash !!!
       buf.insert(insAt, "text to be inserted");
       insAt+=2;
```









Back to the initial problem ...

```
void dolt(const char* in, const char* out)
{
  vector<string> buf;
   vector<string>::iterator insAt = buf.end();
   string linBuf;
   ifstream inFile(in);
   while(getline(inFile,linBuf))
                                    <<<< crash !!!</pre>
         buf.insert(insAt,linBuf);
   sort(buf.begin(), buf.end());
   ofstream outFile(out);
   copy(buf.begin(), buf.end(),
        ostream_i terator<string>(outFile, "\n"));
```

A suggested solution

i nsert() returns an iterator to the newly inserted element.Use this new, valid position as the point of insertion for subsequent insertions.

```
void d(ltconst char* in, const char* out)
{ ...
while(getline(inFile,linBuf))
```

insAt = buf.insert(insAt,linBuf);

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. . .

Another Solution

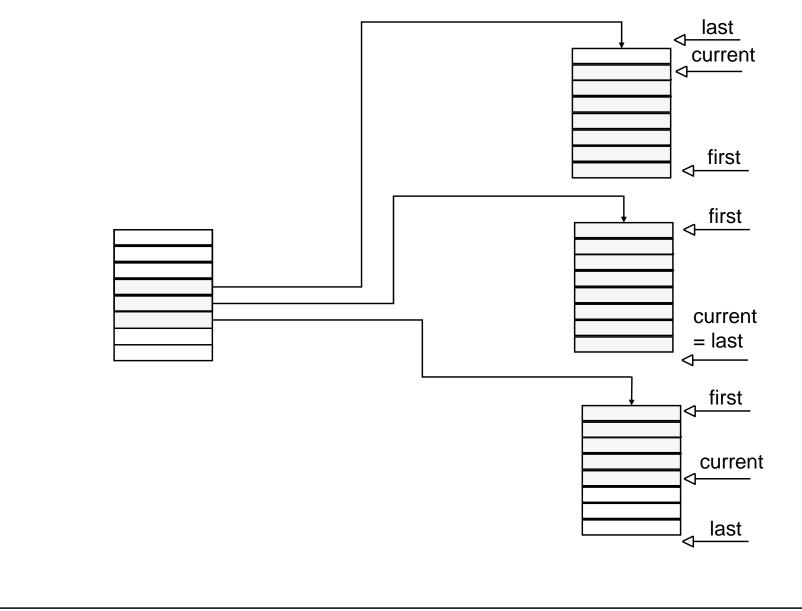
More elegant and easier to comprehend is the use of the push_back() function instead of the i nsert() function.

```
void dolt(const char* in, const char* out)
{
   deque<string> buf;
   string linBuf;
   ifstream inFile(in);
   while(getline(inFile,linBuf))
         buf.push_back(linBuf);
```

Validity guarantee for insertion into a deque:

- An insert in the middle invalidates all the iterators and references.
- An insert at either end of the deque invalidates all the iterators to the deque, but has no effect on the validity of references.

A typical deque implementation



Validity guarantee for erase()

For associative containers and list:

• Invalidates only iterators and references to the erased elements.

For vector:

• Invalidates all the iterators and references after the point of the erase.

For deque:

- An erase in the middle invalidates all the iterators and references to elements of the deque.
- An erase at either end invalidates only the iterators and the references to the erased elements.

STL Pitfall #2

Mutable or Immutable Set Iterators ?

A set container in the STL ...

- is implemented as a binary tree
- needs a strict weak ordering for the elements
- allows insertion and removal of elements via i nsert() and erase(), which use the ordering for maintaining the tree structure
- provides iterators that give access to the contained elements

Note:

- Elements must not be modified through an iterator because direct manipulation of the elements would corrupt the tree structure.
- A set implementation need not provide a mutable iterator.

Conceivable implementations

- Details are still an open issue (#103 on the library issue list of August 1999). Two implementations for set iterators are conceivable:
- constant iterator:
- fool-proof: no chance to modify the elements in-place
- restrictive: cannot change parts of the element that do not contribute to the ordering

mutable iterator:

• security hole: can inadvertently corrupt the tree structure

Consider a program that implements a bank account class, creates a set of bank accounts, and tries to assign to an element through an iterator.

Modification through a set iterator

```
class account {
    ...
    size_t _number; // determines ordering
    double _balance; // irrelevant for ordering
};
```

bool operator<(const account& lhs, const account& rhs)
{ return lhs._number < rhs._number; }</pre>

```
set<account> s;
```

•••

```
set<account>::iterator iter;
```

•••

*iter = *new account; // direct modification of element

Obvious mistake

```
set<account>::iterator iter;
```

*iter = *new account; // direct modification of element

Overwriting an element in the tree structure is likely to destroy the structure.

Result:

...

constant iterator:

– error message; will not compile

mutable iterator:

will corrupt the tree structure; subsequent behavior is unpredictable

```
set<account>::iterator iter;
```

...

*iter = *new account; // direct modification of element

Never "replace" an element in a set; insert the new one and erase the old one.

```
set<account>::iterator iter;
...
s.insert(iter, *new account);
s.erase(iter);
```

Modification through a set iterator

```
class account {
...
size_t _number; // determines ordering
double _balance; // irrelevant for ordering
};
```

```
bool broke(const account& acc)
{ return acc.balance() <= 0; }</pre>
```

```
set<account> s;
...
// remove element if balance is 0 or less
set<account>::iterator garbage;
garbage = remove_if(s.begin(), s.end(), broke), s.end();
s.erase(garbage, s.end());
```

A less obvious mistake

```
set<account>::iterator iter;
```

•••

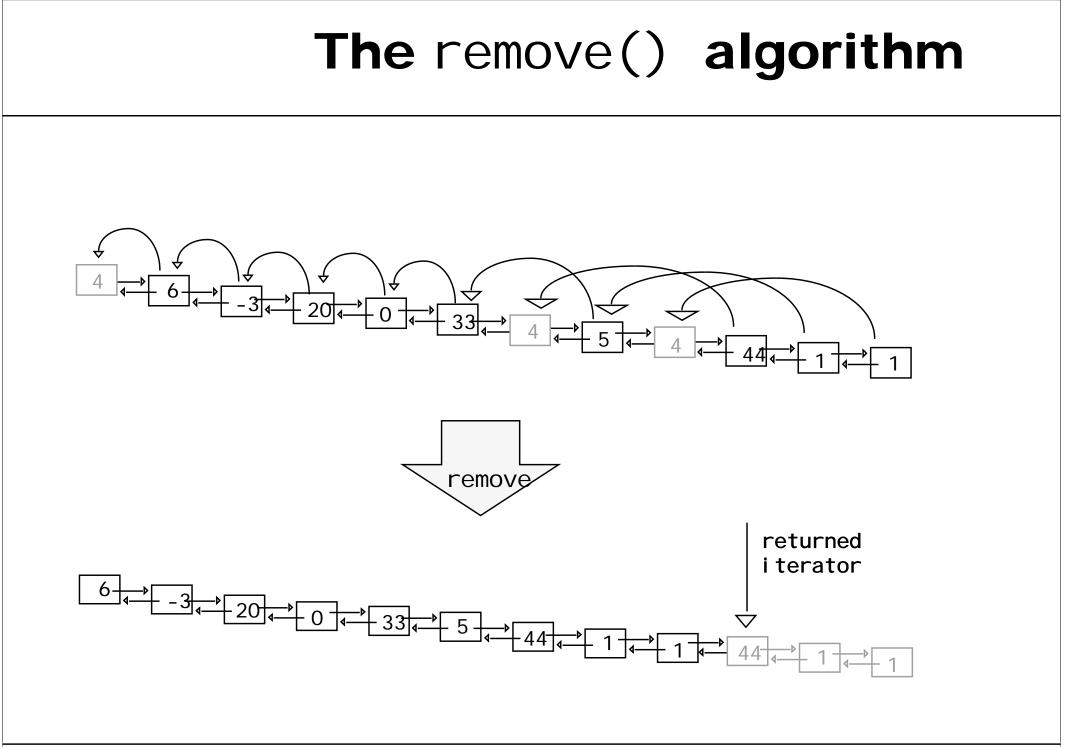
garbage = remove_if(s.begin(), s.end(), broke), s.end();

remove_i f() is a mutating algorithm, that is, it performs inplace modifications on the container elements via the iterator.

Result:

constant iterator:

- error message; will not compile
- warning: "discards const"; will corrupt the tree structure mutable iterator:
- will corrupt the tree structure; subsequent behavior is unpredictable



```
// remove element if balance is 0 or less
set<account>::iterator garbage;
garbage = remove_if(s.begin(), s.end(), broke), s.end();
s.erase(garbage, s.end());
```

Never apply mutating algorithms to a set; instead of (mutating) remove_i f() use inspecting) fi nd_i f().

```
set<account>::iterator fnd;
for (fnd=find_if(s.begin(), s.end(), broke)
;fnd!=s.end()
;fnd=find_if(fnd, s.end(), broke))
{ s.erase(fnd++); }
```

Modifiying part of the element

```
class account {
    ...
    size_t _number; // determines ordering
    double _balance; // irrelevant for ordering
};
bool_operator<(const_account&lbs_const_account)</pre>
```

bool operator<(const account& lhs, const account& rhs)
{ return lhs._number < rhs._number; }</pre>

```
set<account> s;
...
set<account>::iterator iter;
...
// direct modification of part of the element
iter->balance = 1000000;
```

```
set<account>::iterator iter;
```

```
•••
```

// direct modification of part of the element
iter->balance = 1000000;

The balance does not contribute to the ordering relationship. Modifications of the balance would not affect the tree structure.

Result:

constant iterator:

error message; will not compile

mutable iterator:

works nicely

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Conceivable solutions

If the set iterator does not allow modification of the insignificant part of the element:

- Cast away constness.
- Provide a const member function in class account that performs the desired modification.
- Implement an iterator adapter that allows the desired modification.

The Brute Force Approach

```
set<account>::iterator iter;
```

•••

// direct modification of part of the element
iter->balance = 1000000;

Cast away constness:

```
set<account>::iterator iter;
```

•••

// direct modification of part of the element
(const_cast<double>(&(iter->_balance))) = 1000000;

A little more sophisticated

Encapsulate the cast into a const member function of the account class:

```
class account {
public:
 void setBalance(double b) const
 { *const_cast<double*>(&_balance) = b; }
 ...
private:
 size_t _number; // determines ordering
 double _balance; // irrelevant for ordering
};
. . .
iter->setBal ance(1000000);
```

Safety hole: can also modify constant objects of type account

```
set<account>::iterator iter;
```

```
// direct modification of part of the element
iter->balance = 1000000;
```

Define an iterator adapter bal ancel ter that adapts the set iterator. Its dereference operator returns a non-const reference to the balance of the element pointed to.

```
set<account>::iterator iter;
...
// direct modification of part of the element
*balancelter(iter) = 1000000;
```

...

A simple iterator adapter

```
class balancelter {
public:
 explicit balancelter(set<account>::iterator i)
 :_i(i) {}
 double& operator*() const
 { return *const_cast<double*>(&_i ->_balance); }
 balancelter& operator++() { ++_i; return *this; }
 // ... postfix ++, pre- and postfix -- ...
pri vate:
 set<account>::iterator i;
};
```

The 3 suggested solutions

- Casting away constness is the brute force approach; it can and should be avoided.
- Providing a const member function that performs the modification is error-prone; allows modification of const objects.
 - Not a viable solution if the implementation of the account class must not be changed.
- The iterator adapter is the most flexible solution:
 - cast is safely encapsulated;
 - no change to the account class necessary;
 - no security hole; cannot change elements through constant iterators
 - can apply algorithms to the adapted iterator

Conclusions

- It's unfortunate that details of the set iterator are still an open issue.
 - impairs portability efforts
- Avoid direct access through a set iterator to any part that is significant to the ordering of elements.
 - Never allow lvalue use of a dereferenced set iterator.
- Avoid applying mutating algorithms to sets.
 - Read the damned manual.
- If you want to modify a non-significant part and need to get around the const-restriction, build proper abstractions.
 - see for instance the iterator adapter

STL Pitfall #3

Function objects must not have side effects

Function objects in the STL ...

- are accepted as arguments to numerous algorithms
- can be function pointers of functors
- must not have side effects
- must not modify container elements through an iterator

Consider a program that

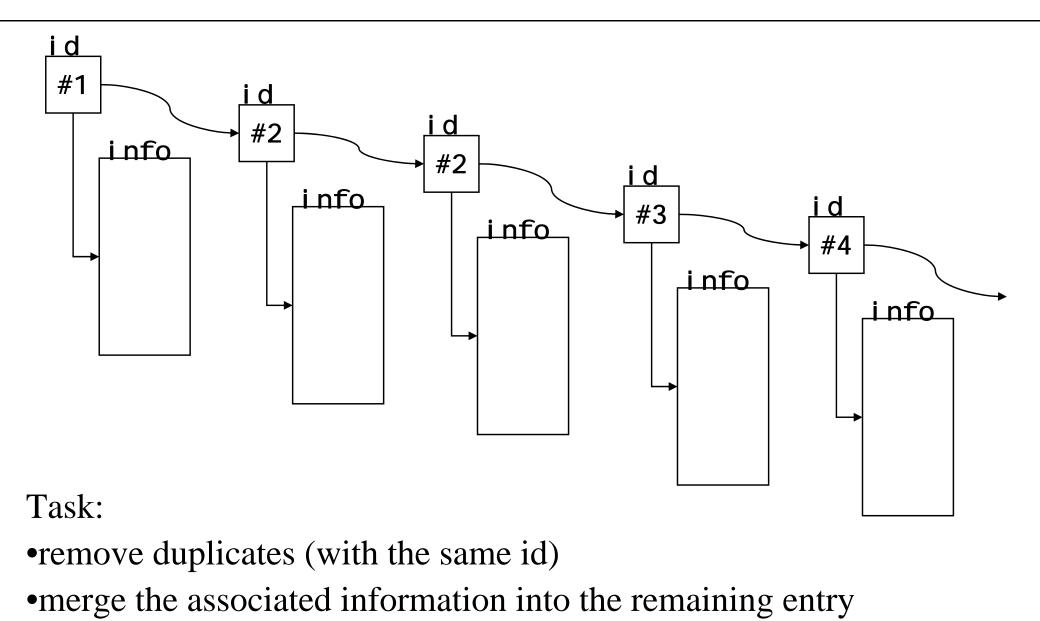
removes duplicates (with the same id) from a container, accumulates any information associated with the id, defines a functor to perform the compression, and calls uni que() passing the functor to get the job done.

An insurance application

```
struct accident {
    string owner; string insurance;
    string date; bool dumped;
};
struct insuranceRec {
    insuranceRec(long id, const list<accident>& c);
    long vehicleld;
    list<accident> crashes;
};
```

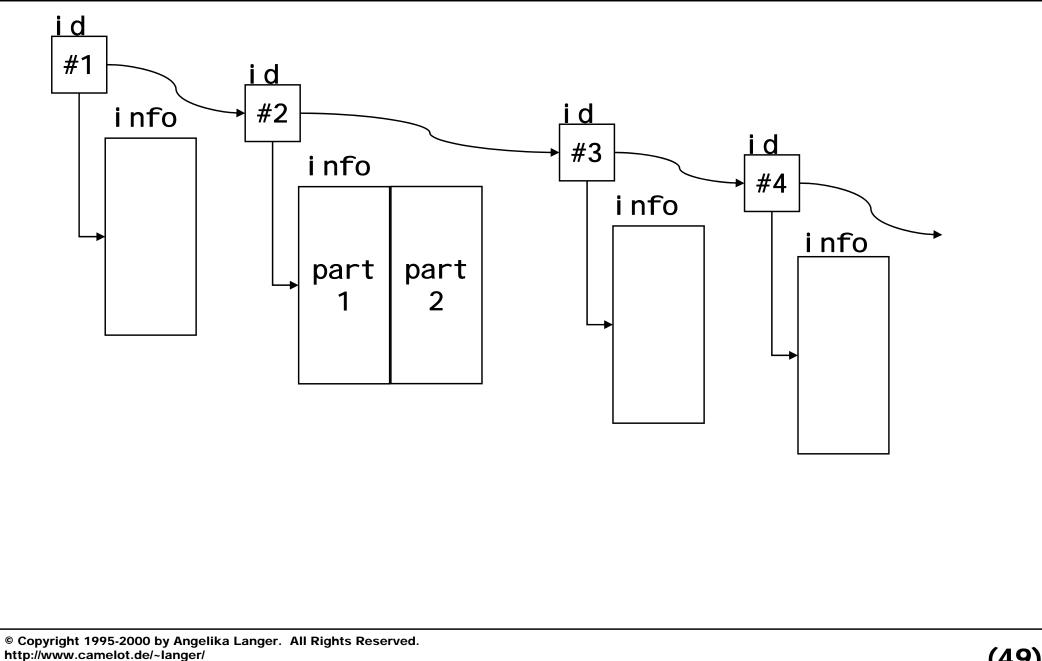
```
multiset<insuranceRec> clients;
// ... populate container ...
clients.erase( // clean up: remove duplicates
    unique(clients.begin(), clients.end(), mergeRec()),
    clients.end());
```

Clean up



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Clean up



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Implementing the clean up

We want to use the uni que() algorithm for elimination of consecutive duplicates:

template < class ForwardI terator, class BinaryPredicate > ForwardI terator unique(ForwardI terator first, ForwardI terator last, BinaryPredicate pred);

We need to define a predicate that

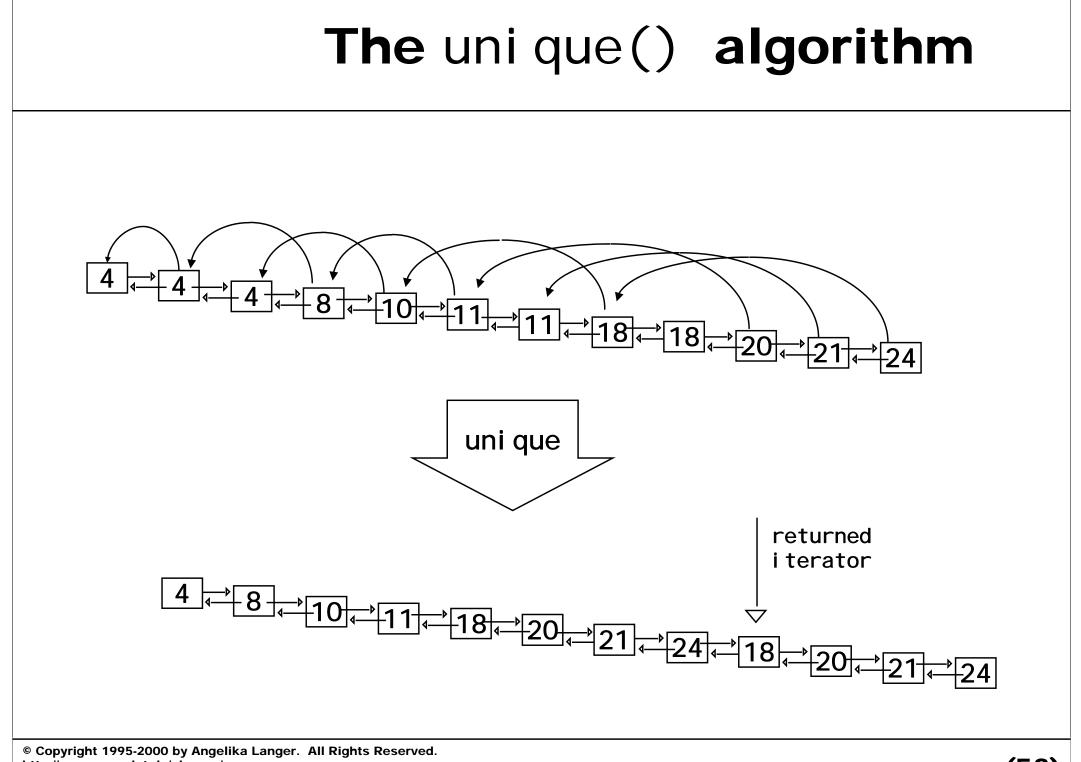
- determines the duplicates (i.e. it must check for identical ids) and
- produces a side effect (i.e. merging the associated info)

The function object type

```
class mergeRec {
public:
   bool operator() ( insuranceRec& lhs,
                    const insuranceRec& rhs)
   { // predicate: check for same id
     bool sameId = (Ihs == rhs);
     if (sameld)
        // produce side effect: append rhs-info to lhs-info
        copy(rhs.crashes.begin(), rhs.crashes.end(),
              inserter(lhs.crashes, lhs.crashes.end());
     return sameld;
```

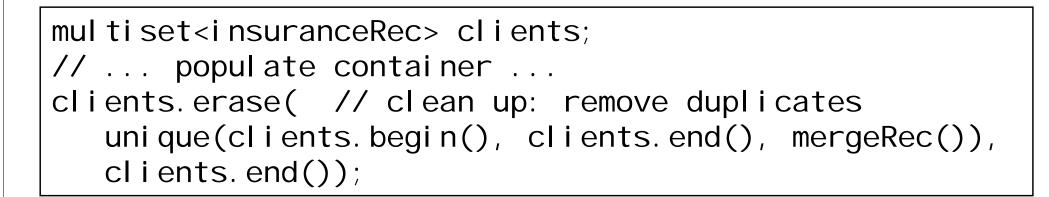
The element type

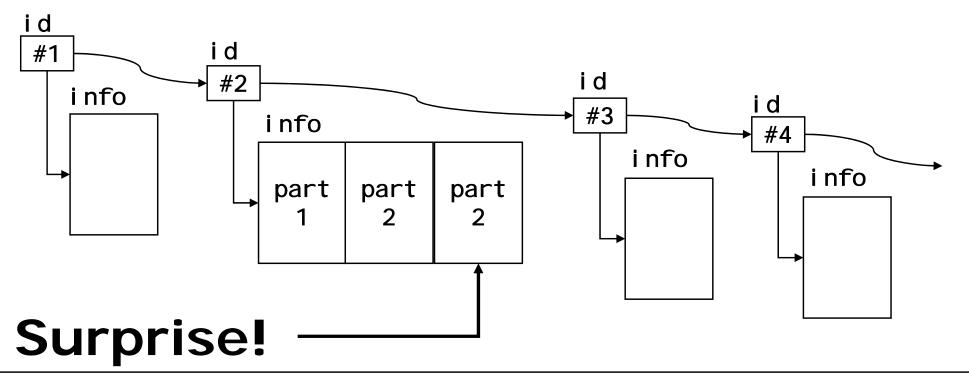
```
struct accident {
   string owner; string insurance;
   string date; bool dumped;
};
struct insuranceRec {
   insuranceRec(long id, const list<accident>& c);
   long vehicleld;
   list<accident> crashes;
};
     operator == (const insuranceRec& lhs,
bool
                const insuranceRec& rhs)
{ return lhs.vehicleId == rhs.vehicleId; }
```



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Doing the clean up





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What happened?

• Our predicate does not only check for the duplicates but in addition produces a side effect (i.e. merging the associated info).

Function objects must not have side effects and must not modify any element through an iterator.

• Both is violates, which does not always create problems, but in our case we do not know how often the uni que() algorithm produces the predicate's side effect.

It is unspecified how often an algorithm invokes a function object on the same (pair of) element(s).

Implementation of unique()

```
template <class ForwardIterator, class BinaryPredicate>
ForwardIterator
unique(ForwardIterator first, ForwardIterator last,
            BinaryPredicate binary_pred)
{
    first = adjacent_find(first, last, binary_pred);
    return unique_copy(first, last, first, binary_pred);
```

adj acent_find() return the first element of a series of duplicates. uni que_copy() copies all elements except for consecutive duplicates.

}

Conceivable solutions

- Do not provide functions (or function objects) that produce side effects to any STL algorithm.
 - implement side effects separately
 - use STL algorithms only for side-effect-free operations
- Implement your own version of unique() so that you have control over the number of side effects produced.

Side-effect-free predicate

Instead of calling uni que() with our mergeRec predicate

- call adj acent_find() with a side-effect-free predicate eqRec and equal_range() to identify any duplicates and
- produce the side effect independently of the STL.

The actual logic

```
multiset<insuranceRec> clients;
// ... populate container ...
typedef multiset<insuranceRec>::iterator iterType;
iterType duplicate = clients.begin();
while (duplicate!=clients.end())
  // identify duplicates
{
   duplicate =
      adj acent_find(duplicate, clients.end(), eqRec());
   // merge info and erase duplicates
   if (duplicate! = clients. end())
   { pair<iterType, iterType> range;
     range = clients.equal_range(*duplicate);
     compress(range);
     clients.erase(++(range.first), range.second);
```

Producing the side-effect

```
template <class Iterator>
void compress(pair<Iterator, Iterator> range)
{
    range. second--;
    copy(range. second->crashes. begin(),
        range. second->crashes. end(),
        inserter(range. first->crashes,
            range. first->crashes. end())
    );
}
```

Customizing unique()

Instead of calling the STL uni que() implement your own version in order to gain control of the side effects produced.

- pro:
 - portable, reusable
 - potentially more efficient
- con:
 - extra effort required

User-defined version of unique()

```
template <class ForwardIterator, class BinaryPredicate>
Forwardl terator
my_unique(ForwardIterator first, ForwardIterator last,
          BinaryPredicate binary_pred)
{ if (first == last) return last;
  else { ForwardIterator next = first;
         while(++next != last)
         { if (!binary_pred(*first, *next)) first = next;
           else // duplicate found
           { while (++next != last)
              if (!binary_pred(*first, *next))
                 *++first = *next;
              return ++first;
         return last;
```

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Evaluation

- using STL version of uni que() with a side-effect producing predicate
 - elegant and readable, but non-portable
- using customized version of uni que() with a side-effect producing predicate
 - still elegant and readable, but also portable
 - extra effort required
- strict separation between side-effect-free predicate and a side-effect producing function
 - comparable to customized unique() regarding effort and complexity
 - portable, but probably not reusable

Further problematic cases

- function objects that modify any part of an element that is relevant to the ordering
 - There is no mechanism to make sure that only immutable iterators are provided to a function object.
 - Hence you can change whatever you like even corrupt the container.
- function objects that depend on how often they are invoked
 - The number of produced side effects is not specified.
 - Any side-effect-producing function object falls into this category.
- function objects that depend on how often they are copied, assigned, or destroyed
 - It is unspecified how many temporary copies of a function object ara algorithms creates.
 - All function objects that have non-constant state and accumulate data between subsequent invocations fall into this category.

Function objects with state

Example:

• Our side-effect producing predicate might make a note (in an internal list) of all duplicates that it finds.

Problem:

- Hardly any algorithm returns the function objects that it received.
 - How do I get hold of the accumulated data?
 - The predicate writes the entire information to a global or static location when it is destroyed.

References to function objects

for_each() is the only algorithm that returns the function
 object.

An alternative approach for getting access to the object state:

- Pass the function objects by reference rather than by value.
 - Requires explicit function arguments specification syntax and creates lifetime dependencies.

unique(clients.begin(), clients.end(), mergeRec());

would become

Destructor with side effects

- The predicate writes the entire information to a global or static location when it is destroyed.
- The destructor has a side effect.
- It is unspecified
 - how many temporary copies of the function object are created inside uni que() and
 - how often the destructor is invoked.

Rule:

- Never create functions objects that produce side effects when they are created, copied, or destroyed.
- This is common sense for any class, but even more important for types that are provided to the STL.

Implementation of unique()

```
template <class ForwardIterator, class BinaryPredicate>
ForwardIterator
unique(ForwardIterator first, ForwardIterator last,
            BinaryPredicate binary_pred)
{
    first = adjacent_find(first, last, binary_pred);
    return unique_copy(first, last, first, binary_pred);
```

We cannot tell how often the predicate is copied unless we also study the implementations of adj acent_find() and uni que_copy().

Restrictions to function objects

- Never create functions objects that produce side effects when they are created, copied, or destroyed.
 - This is common sense for any class, but even more important for objects that are provided to the STL.
- Be careful with function objects that create side effect when they are invoked.
 - It is not at all uncommon that functions have side effects.
- Be careful with function objects that modify the elements through the iterator.
 - It is not at all uncommon that functions which take pointers/references modify the pointed to objects, and iterators have pointer-like semantics.

STL Pitfall #4

Comparators must not be polymorphic

Polymorphic comparators

- It's a common technique to implement a family of compare policies as a class hierarchy with a common base class and to invoke the policies through base class reference for polymorphic behavior. (See the Strategy pattern à la GOF.)
- The associative containers accept a reference to a comparator object; hence one could pass a base class reference.
- However, they store a copy of the comparator object internally, the result of which is object slicing.

STL Pitfall #5

Equality vs. induced equivalence

Equality vs. equivalence

- The associative containers use an induced equivalence relation for maintaining the underlying tree structure.
- The equivalence is educed from the ordering, i.e., the comparator that is provided to the container.
- Container member functions use the equivalence relations for finding equal elements in the container.
- STL algorithms use an equality relation on the iterator's value type for finding equal elements.

Equality vs. equivalence

• Can lead to surprising results when equality and equivalence are different.

Example: case-insensitive string compare

- Strings that are equivalent (regarding case) are not necessary equal.
- Consider a multiset of strings with case-insensitive order in conjunction with "set" algorithms such as uni on() or intersection(.

STL Pitfall #6

Type incompatibility of adapted iterators

Adapted iterators

- STL algorithms accept any kind of iterator type, because they are function templates.
- Container member functions only accept their own iterator type.

Example:

- A container element is searched for by passing reverse iterators to find().
- The resulting iterator cannot be passed to the container's erase() function.
- When implementing an iterator adapter, never forget ot implement the base() member function.

STL Pitfall #7

Several stream iterators on the same stream

Interdependent stream iterators

- Stream iterators on the same stream are not independent of each other.
- Advancing one iterator affects all other iterators on the same stream, because it changes the underlying stream position.
- For input stream iterators:
 - Increment means reading from the stream.
 - Dereferencing means providing the stored, previously read value.
- For output stream iterators:
 - Increment and dereferencing are NOPs.
 - Assignment to the iterator means writing to the stream.

Interdependent stream iterators

- Reaching the end iterator means failure of the read/write operation, i.e. reaching end of input or an error situation.
- Example: 2 input stream iterators for reading a string and a float
- Reading a string also moves the float iterator.
- Reading a float when there is a string on the file turns the float iterators into an end iterator; the string iterator can still read.
 - The frozen float iterator cannot be reset.
 - The stream state must be cleared and a new float iterator must be created.

STL Pitfall #8

Allocators must exhibit static behavior

Static allocators

Allocators of the same type must be interchangeable, i.e. must not have state.

Reason:

- Two containers of the same type, i.e. using the same type of allocator, can have different allocator objects.
 - Example: database allocators to different databases
- If the two containers are assigned to each other, all elements must be copied. Which allocator must be used?
- There's no universal answer.
- The problem evaporates when all allocators of the same exhibit the same behavior.

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